

WHAT IS CLAIMED IS:

1. A code preparation method comprising:
 identifying at least one operation in first executable instance of code;
 executing the first executable instance and responsive to detection of an
 execution event, associating a corresponding execution characteristic
 with a corresponding identified one of the operations; and
 preparing a second executable instance of the code based, at least in part, on
 the association between the execution characteristic and the identified
 operation.
2. The method of claim 1,
 wherein the operation identification is consistent between the first executable
 instance and the preparation of the second executable instance.
3. The method of claim 2,
 wherein the consistency of operation identification is maintained from
 preparation of the first executable instance to preparation of the second
 executable instance.
4. The method of claim 1,
 wherein same unique identification numbers are assigned to corresponding
 operations of the first executable and the second executable.
5. The method of claim 4,
 wherein the execution characteristic is associated with the unique
 identification number.
6. The method of claim 4,
 wherein the unique identification numbers and their assignment to operations
 are maintained throughout any optimizations or code transformations
 performed in preparation of the first executable.

7. The method of claim 6,
wherein the maintenance of the unique identification number assignments
include further assigning the unique identification number to a copy
when an operation is copied as part of a code transformation or
optimization.
8. The method of claim 6,
wherein the maintenance of the unique identification number assignments
includes removing an assignment when the assigned operation is
removed as part of a code transformation or optimization.
9. The method of claim 1,
wherein the associating of the corresponding execution characteristic includes
encoding aggregated hardware event information in an extended
definition of an instruction instance for use in the preparation of the
second executable instance.
10. The method of claim 1,
wherein the identified operation is a memory access instruction.
11. The method of claim 1,
wherein the execution characteristic includes a cache miss likelihood.
12. The method of claim 1,
wherein the preparation includes inserting one or more prefetch operations in
the code prior to the identified operation to exploit latency provided by
servicing of a cache miss by the identified operation.
13. The method of claim 1, further comprising:
preparing the first executable instance.
14. The method of claim 13,
wherein the preparation of the first executable instance includes substantially
all optimizations operative in the preparation of the second executable.

15. The method of claim 14,
wherein execution of the first executable instance corresponds substantially
with execution of an executable instance of code prepared without the
identifying.
16. The method of claim 14,
whereby execution of the first executable instance sufficiently corresponds to
that in an expected execution environment, so that the execution
characteristic is applicable to the preparation of the second executable.
17. The method of claim 13,
wherein the preparation of the first executable instance forgoes certain
optimizations performed, after use of the association between the
execution characteristic and the identified instruction, by the further
preparing.
18. The method of claim 13,
wherein the preparation of the first executable instance includes compilation
of the code.
19. The method of claim 1,
wherein both the first and the second executable instances are compiled
instances of the code.
20. The method of claim 1,
wherein the second executable instance is an optimization of the first
executable instance.
21. The method of claim 1,
wherein the preparing includes optimizations forgone in the first executable
instance.

22. The method of claim 1,
wherein the preparation of the second executable instance includes
optimizations forgone in preparation of the first executable instance.
23. The method of claim 1,
wherein at least the preparing is performed by an optimizing compiler.
24. The method of claim 1,
wherein at least the preparing is performed by a binary translator.
25. The method of claim 1,
wherein at least the preparing is performed by a binary rewriter.
26. The method of claim 1,
wherein at least the preparing is performed by a binary optimizer.
27. The method of claim 1,
wherein at least the preparing is performed by a just-in-time (JIT) compiler.
28. The method of claim 1,
wherein the associating of the corresponding execution characteristic includes
aggregating contributions of plural instances of the execution event.
29. The method of claim 1,
wherein the associating of the corresponding execution characteristic includes
backtracking from a point in the code that coincides with delayed
detection of the execution event.
30. The method of claim 1,
wherein the associating of the corresponding identified one of the operations
includes reading or receiving a computer readable encoding of an event
profile.

31. The method of claim 1,
wherein the associating of the corresponding execution characteristic includes
reading or receiving a computer readable encoding of an event profile.

32. The method of claim 1, further comprising:
preparing the second executable instance as a computer program product for
distribution, transmission or execution.

33. The method of claim 33,
wherein the computer program product is encoded in one or more computer
readable media selected from the set of a disk, tape or other magnetic,
optical, semiconductor or electronic storage medium and a network,
wireline, wireless or other communications medium.

34. An optimizing compiler that prepares a second executable instance of
computer program code including optimizations in addition to those of a previously
prepared first executable instance thereof, wherein the additional optimizations
include performing one or more transformations based on run-time information from
execution of the first executable instance, wherein consistency of instruction
identification is maintained from preparation of the first executable instance to
preparation of the second executable instance.

35. The method of claim 34,
wherein same unique identification numbers are assigned to corresponding
operations of the first executable and the second executable.

36. The method of claim 35,
wherein the unique identification numbers and their assignment to operations
are maintained throughout any optimizations or code transformations
performed in preparation of the first executable.

37. The method of claim 36,
wherein the maintenance of the unique identification number assignments
include further assigning the unique identification number to a copy

when an operation is copied as part of a code transformation or optimization.

38. The method of claim 36,
wherein the maintenance of the unique identification number assignments
includes removing an assignment when the assigned operation is
removed as part of a code transformation or optimization.

39. The optimizing compiler of claim 34,
wherein the transformations include insertion of one or more prefetch
instructions.

40. The optimizing compiler of claim 34,
wherein the transformations include insertion of one or more non-faulting
loads.

41. The optimizing compiler of claim 34,
wherein selection of optimizations performing in the preparation of the first
executable instance is biased toward collection of data.

42. The optimizing compiler of claim 34,
wherein the additional optimizations performing in the preparation of the
second executable instance are biased toward obtaining improved
performance based on the run-time information.

43. The optimizing compiler of claim 34,
wherein transformations include insertion of instructions into the second
executable instance to reduce latency of memory access operations
that, based on the run-time information, are likely to miss in a cache.

44. The optimizing compiler of claim 34,
wherein the optimizing compiler prepares the second executable instance, but
not the first.

45. The optimizing compiler of claim 34,
wherein the optimizing compiler also prepares the first executable instance of
computer program code.

46. The optimizing compiler of claim 34, embodied as part of a binary
translator.

47. The optimizing compiler of claim 34, embodied as part of a binary
rewriter.

48. The optimizing compiler of claim 34, embodied as part of a binary
optimizer.

49. The optimizing compiler of claim 34, embodied as a just-in-time (JIT)
compiler.

50. The optimizing compiler of claim 34,
wherein first and second executions of the optimizing compiler respectively
provide the first and second executable instances; and
wherein the transformations are performed in addition to optimizations
coextensive with those performed in the first executable instance.

51. The optimizing compiler of claim 34,
wherein the optimizing compiler identifies one or more memory access
instructions in the first executable instance of the computer program
code; and
wherein the run-time information encodes respective execution characteristics
for respective ones of the identified memory access instructions.

52. The optimizing compiler of claim 34,
wherein collection of the run-time information includes aggregation of
execution event information and association of the aggregated
information with memory access instructions identified in the first
executable instance of the computer program code.

53. The optimizing compiler of claim 34,
encoded in one or more computer readable media selected from the set of a
disk, tape or other magnetic, optical, semiconductor or electronic
storage medium and a network, wireline, wireless or other
communications medium.

54. A method of optimizing code for an execution environment in which a
possibility of processor or pipeline stall latency exists for particular instructions
thereof, the method comprising:

identifying the particular instructions in a first executable instance of the code;
associating a characterization of stall likelihood with respective ones of the
particular instructions based on at least one execution of the first
executable instance; and
inserting behind respective ones of the particular instructions, one or more pre-
executable portions of the particular instructions selected to reduce
stall latency thereof based on the respective associated characterization
of stall likelihood.

55. The method of claim 54,
wherein the pre-executable portions include prefetch instructions.

56. The method of claim 54,
wherein the pre-executable portions include non-faulting loads.

57. The method of claim 54,
wherein the particular instructions are memory access instructions and the
associated characterizations are of cache miss likelihood.

58. The method of claim 54, further comprising:
executing the first executable instance of the code to obtain the
characterization of stall likelihood for the particular instructions.

59. The method of claim 54, further comprising:
preparing a computer program product encoding a second executable instance
of the code that includes the inserted prefetch instructions.

60. The method of claim 54, further comprising:
preparing the first executable instance of the code.

61. A computer program product encoded in one or more computer readable
media, the computer program product comprising:
a first execution sequence; and
an information encoding associating an execution event with at least some
operation of the first execution sequence, the associated execution
event based at least in part on an execution profile of the first
execution sequence of operations, wherein consistency of the
association is maintained from preparation of the first executable
instance for preparation of a second executable instance.

62. The computer program product of claim 61,
wherein the execution event is a cache miss likelihood.

63. The computer program product of claim 61,
wherein the associated operation is a memory access operation.

64. The computer program product of claim 61,
employed in an data structure of an optimizing compiler in preparation of an
optimized instance of the execution sequence of operations, wherein
the optimized instance includes one or more prefetch operations placed
before particular ones of the memory access operations for which the
associated information encoding indicates a cache miss likelihood.

65. The computer program product of claim 61,
wherein the one or more computer readable media are selected from the set of
a disk, tape or other magnetic, optical, semiconductor or electronic

storage medium and a network, wireline, wireless or other communications medium.

66. An apparatus comprising:

means for identifying in a first executable instance of computer program code

certain operations thereof for inclusion in an execution profile;

means for collecting the execution profile; and

means for preparing a second executable instance of the computer program

code, wherein consistency of the identifying is maintained for

operations thereof corresponding to the certain operations such that the

corresponding certain operations are relatable to the execution profile.

67. The apparatus of claim 66,

wherein the identifying includes producing a table of tags and operation addresses.

68. The apparatus of claim 66,

wherein information for the identifying is encoded in a file or communications channel read by the means for collecting.

69. The apparatus of claim 66, further comprising:

means for preparing the first executable instance of the computer program code.